

# Social Networks and Belief Merge Communication Protocols

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Beijing, October 18 2013

## Outline

### 1) Seligman, Girard & Liu (2011, 2013)

- ▶ social network
- ▶ peer pressure effects,  
influence inbetween  
“friends”



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- ▶ plausibility
- ▶ effects of group members sharing information with the rest of the group



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### 2) Baltag & Smets (2009)

- ▶ plausibility
- ▶ effects of group members sharing information with the rest of the group



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### 3) Aim: a unified social network plausibility framework

- ▶ model social influence on beliefs through communication among agents in a social network
- ▶ define some particular communication protocols (in the new framework) inspired by 2) to represent some level of influence as defined in 1)



## 1) Social influence à la Girard, Liu & Seligman



### The framework

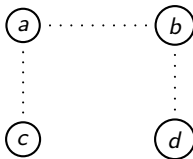
Static hybrid logic to represent who is friend with whom and who believes what  
+ an (external) influence operator

### The main ideas

- ▶ Agents are influenced by their friends and only by their friends.
- ▶ Simple “peer pressure principle”: I tend to align with my friends.
- ▶ “Being influenced” is defined as “aligning my beliefs to the ones of my friends”.
- ▶ No communication is (at least explicitly) involved. (transparency?)

## Friends network

Social network frame:



- ▶ *a* is friend with agents *b* and *c*
- ▶ *b* is *d*'s only friend
- ▶ *a* is *c*'s only friend.

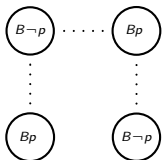
## Belief revision induced by (direct) social influence

### 3 possible states

- ▶  $Bp$
- ▶  $B\neg p$
- ▶  $Up := \neg Bp$  and  $\neg B\neg p$

### Strong influence

When all of my friends believe that  $p$ , I (successfully) *revise* with  $p$ . When all of my friends believe that  $\neg p$ , I (successfully) *revise* with  $\neg p$ .



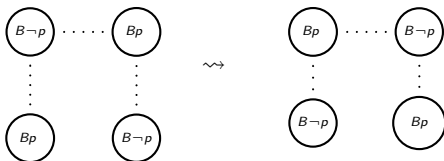
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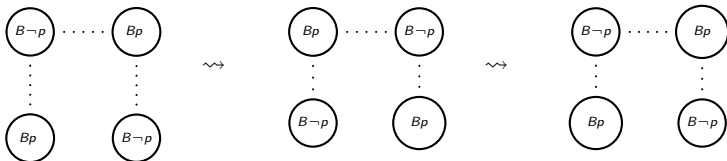
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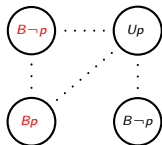
## Belief contraction induced by social influence

### Weak influence

None of my friends supports my belief in  $p$  and some believe that  $\neg p$ .

I (successfully) *contract* it.

(And similarly for  $\neg p$ )



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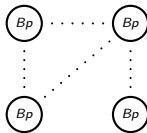
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(And similarly for  $\neg p$ )



## Stabilization

- ▶ Stable state: applying the social influence operator doesn't change the state of any agent.
- ▶ Stabilization: some configurations will reach a stable state after a finite number of applications of the influence operator (see example of weak influence above) and some won't (see example of strong influence).
- ▶ Sufficient condition for stability: all friends are in the same state.



## 2) Communication protocols à la Baltag & Smets



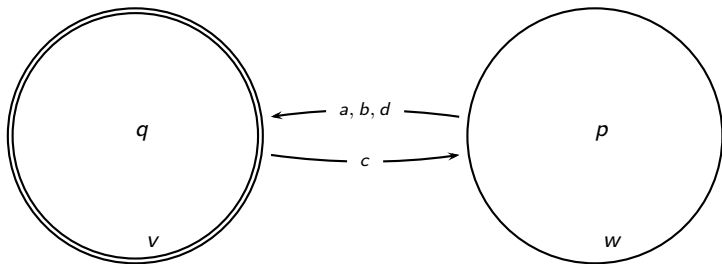
### The framework

DEL type: plausibility modeling of (several) doxastic attitudes + communication events

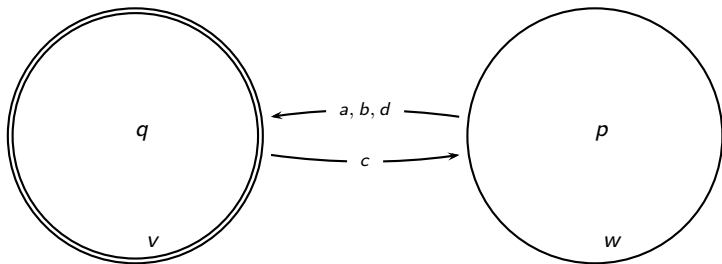
### The main ideas

- ▶ Agents communicate via public announcements.
- ▶ Assuming that they trust each other enough, agents all revise their beliefs with each of the announced formula, sequentially.
- ▶ In this sense, each announcement influences everybody (else) into belief revision.

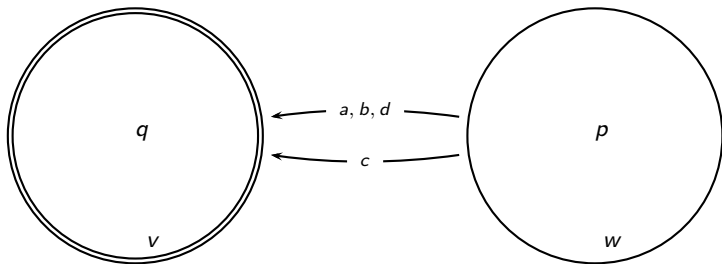
## Plausibility model



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## Reaching a stable state of agreement

### How to communicate?

- ▶ Agents speak in turn (given expertise rank).
- ▶ An agent announces all and only (non-equivalent) sentences that she believes (honesty + exhaustivity).
- ▶ After a finite number of announcements (and corresponding revisions), everybody holds the same beliefs.
- ▶ This is a stable state: nothing which could be announced by any agent would change anything anymore.

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### Lexicographic belief merge protocol

$$\rho_a := \prod \{ \uparrow \phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}, w \models B_a \phi \}$$

$$\rho_b := \prod \{ \uparrow \phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}_{[\rho_a]}, w \models B_b \phi \}$$

etc for all  $c \in \mathcal{A}$

where  $\prod$  is a sequential composition operator and  $\mathcal{M}_{[\rho_a]}$  is the new model after joint revision with each formula announced by  $a$ .

## Big picture

### Common features

- ▶ Agents are influenced into revising their beliefs to make them closer to the ones of (some) others.
- ▶ A global agreement state is stable (both under honest communication and under social conformity pressure).

From 1)



- ▶ **Social network**
- ▶ Synchronic
- ▶ **Over friends only**
- ▶ Equal power (among friends)
- ▶ Direct
- ▶ **Tools:** nominals, @,  $F$

From 2)

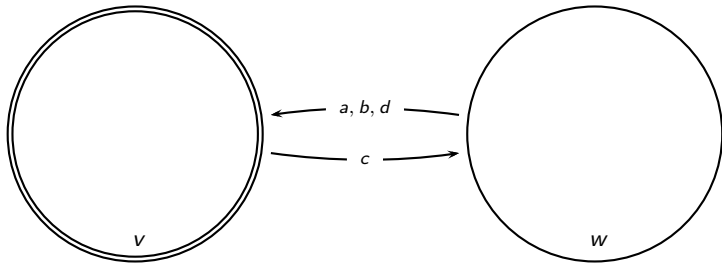


- ▶ **Plausibility**
- ▶ Sequential
- ▶ Over everybody
- ▶ Ranking
- ▶ **Via communication**
- ▶ **Tools:**  $B, \uparrow, \uparrow$

### 3) A social network plausibility framework

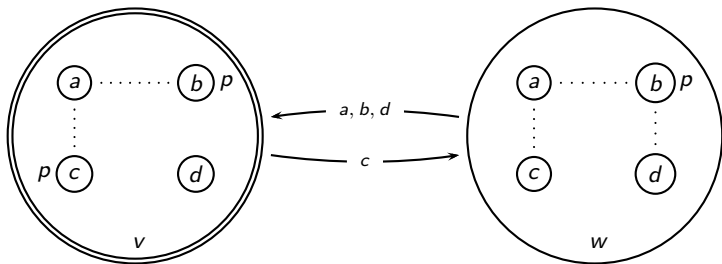


plausibility model:



### 3) A social network plausibility framework +

Social network plausibility model:



## Social network plausibility model

$$\mathcal{M} = (S, \mathcal{A}, \leq_a \subseteq S \times S, \|\cdot\|, s_0, \succ_s \subseteq S \times S)$$

- ▶  $S$  is a (finite) set of possible states.
- ▶  $\mathcal{A}$  is a (finite) set of agents.
- ▶  $\leq_a \subseteq S \times S$  is a locally connected preorder, interpreted as the subjective plausibility relation of agent  $a$ , for each  $a \in \mathcal{A}$
- ▶  $s_0 \in S$  is a designated state, interpreted as the actual state
- ▶  $\succ_s \subseteq S \times S$  is an irreflexive and symmetric relation, interpreted as friendship, for each state  $s \in S$
- ▶  $\|\cdot\| : \Phi \cup N \rightarrow \mathcal{P}(S \times \mathcal{A})$  is a valuation, assigning:
  - ▶ a set  $\|p\| \subseteq S \times \mathcal{A}$  to every element  $p$  of some given set  $\Phi$  of “atomic propositions”
  - ▶ a set  $\|n\| = S \times \{a\}$  for some  $a \in \mathcal{A}$  to every element  $n$  of some given set  $N$  of “nominals”.

## Syntax

$$\phi := p \mid n \mid \neg\phi \mid \phi \wedge \phi \mid F\phi \mid @n\phi \mid B\phi$$

where  $p$  belongs to a set of atomic propositions  $\Phi$  and  $n$  to a set of nominals  $N$ .

## Inherited indexicality

Formulas evaluated both at a state  $w \in S$  and at an agent  $a \in A$ .

- ▶  $p$  : "I have a moustache."
- ▶  $BFp$ : "I believe that all my friends have a moustache."
- ▶  $FBp$ : "All of my friends believe that they have a moustache".



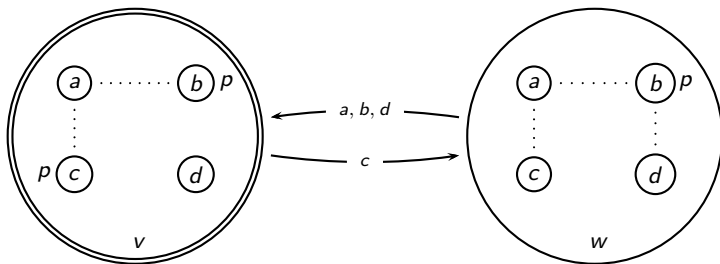
## Semantic clauses

- ▶  $\mathcal{M}, w, a \models p$  iff  $\langle w, a \rangle \in \llbracket p \rrbracket$
- ▶  $\mathcal{M}, w, a \models n$  iff  $\langle w, a \rangle \in \llbracket n \rrbracket$  iff  $a = \underline{n}$
- ▶  $\mathcal{M}, w, a \models \neg\phi$  iff  $\mathcal{M}, w, a \not\models \phi$
- ▶  $\mathcal{M}, w, a \models \phi \wedge \psi$  iff  $\mathcal{M}, w, a \models \phi$  and  $\mathcal{M}, w, a \models \psi$
- ▶  $\mathcal{M}, w, a \models F\phi$  iff  $\mathcal{M}, w, b \models \phi$  for all  $b$  such that  $a \succ b$
- ▶  $\mathcal{M}, w, a \models @b\phi$  iff  $\mathcal{M}, w, \underline{b} \models \phi$
- ▶  $\mathcal{M}, w, a \models B\phi$  iff  $\mathcal{M}, v, a \models \phi$  for all  $v \in S$  such that  $v \in best_a w(a)$

notation:

- ▶  $\underline{n}$  the unique agent at which the nominal  $n$  holds
- ▶  $s(a)$  the comparability class of state  $s$  relative to agent  $a$ :  $t \in s(a)$  iff  $s \leq_a t$  or  $t \leq_a s$
- ▶  $best_a s(a)$  the most plausible states in  $s(a)$  according to  $a$ :  $best_a s(a) := \{s \in s(a) : t \leq_a s \text{ for all } t \in s(a)\}$

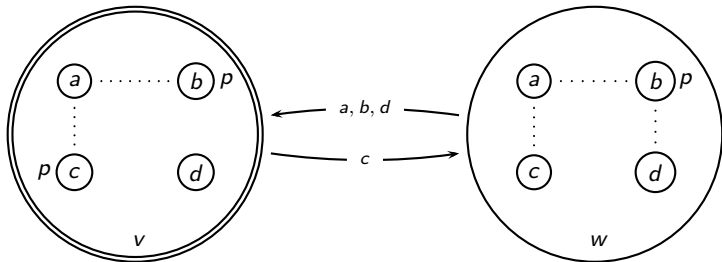
## Example



- ▶  $M, v, \underline{c} \models p$
- ▶  $M, v, \underline{a} \models Fp$
- ▶  $M, v, \underline{a} \models \langle F \rangle b$

- ▶  $M, w, \underline{d} \models FBp$
- ▶  $M, w, \underline{a} \models BFp$


## Example

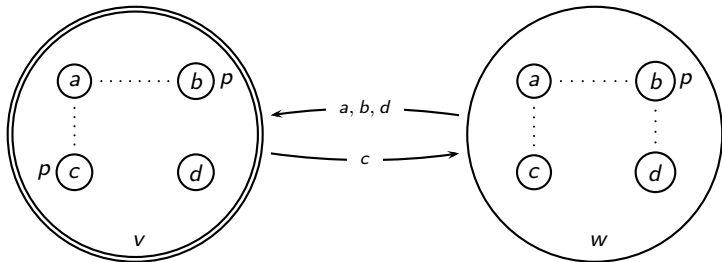


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$\neg @b \langle F \rangle d$  

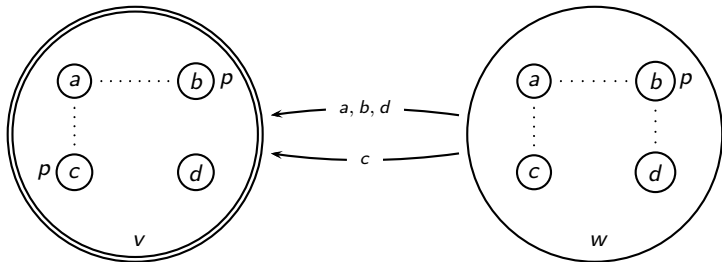


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$\neg @b \langle F \rangle d$  



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## Influence dynamics

### Simplifying assumptions

- ▶ agents speak in turn (rank)
- ▶ only friends communicate
- ▶ agents revise with (all) sentences announced (trust)

## Revision operator

### Joint radical upgrade $\uparrow \phi$

- ▶ “Promote” all the  $\|\phi\|$ -worlds so that they become more plausible than all  $\neg\|\phi\|$ -worlds (in the same information cell), keeping everything else the same:

## Revision operator

### Joint radical upgrade $\uparrow \phi$

- ▶ “Promote” all the  $\|\phi\|$ -worlds so that they become more plausible than all  $\neg\|\phi\|$ -worlds (in the same information cell), keeping everything else the same:
- ▶  $\uparrow \phi$  is a model transformer which takes as input any model  $\mathcal{M} = (S, \mathcal{A}, \leq_{a \in \mathcal{A}}, \|\cdot\|, s_0, \succ_{s \in S})$  and outputs a new model  $\mathcal{M}' = (S, \mathcal{A}, \leq'_{a \in \mathcal{A}}, \|\cdot\|, s_0, \succ_{s \in S})$  such that:  
 $s \leq'_a t$  iff either  $(s, t \notin \|\phi\| \text{ and } s \leq_a t)$  or  $(s, t \in \|\phi\| \text{ and } s \leq_a t)$  or  $(t \in s(a) \text{ and } s \notin \|\phi\| \text{ and } t \in \|\phi\|)$ .



## Belief merge

Baltag & Smets' lexicographic belief merge protocol

$$\rho_a := \prod \{ \uparrow \phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}, w \models B\phi \}$$

$$\rho_b := \prod \{ \uparrow \phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}_{[\rho_a]}, w \models B\phi \}$$

etc for all  $c \in \mathcal{A}$

where  $\prod$  is a sequential composition operator and  $\mathcal{M}_{[\rho_a]}$  is the new model after joint revision with each formula announced by  $a$ .

## Belief merge

### Indexical lexicographic belief merge protocol

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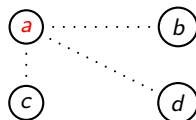
etc for all  $c \in \mathcal{A}$

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## A central friend

### Assumptions

- ▶  $a$  is other agents' only friend.
- ▶  $a$  speaks first.



### One-to-others unilateral strong influence protocol

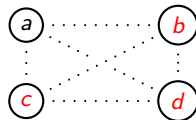
One step version of the indexical lexicographic belief merge protocol:

$$\rho_a := \prod \{ \uparrow @_a \phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}, w, \underline{a} \models B\phi \}$$

## Everybody is friends with everybody else

### Assumption

- ▶ Connectedness



### Others-to-one unilateral strong influence protocol

$$\rho_b := \prod \{ \uparrow @_b B\phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}, w, \underline{b} \models B\phi \}$$

$$\rho_c := \prod \{ \uparrow @_c B\phi : \|\phi\| \subseteq S \times \mathcal{A} \text{ such that } \mathcal{M}, w, \underline{c} \models B\phi \}$$

etc, for all  $d \in \mathcal{A}$  such that  $\mathcal{M}, w, d \models \langle F \rangle a$

$$\rho_a := \prod \{ \uparrow @_a \phi \text{ iff } \mathcal{M}_{[\rho_b; \rho_c; \dots]}, w, \underline{a} \models BFB\phi \}$$

where  $\mathcal{M}_{[\rho_b; \rho_c; \dots]}$  is the model resulting from the successive revisions (by all friends) with each of the formulas announced by each of them.

## Summary

- ▶ Social network plausibility framework with communication events
- ▶ Indexical protocol to merge beliefs
- ▶ Unilateral strong influence *one-to-all-the-others* protocol
- ▶ Unilateral strong influence *all-the-others-to-one* protocol


## To do next

- ▶ Private (and synchronic?) communication: *friends to friends* influence (level of privacy to determine)
- ▶ Different doxastic attitudes (conditional belief, strong belief, safe belief) + different levels of trust (dynamic attitudes) corresponding to different types of revision (minimal revision, update).
- ▶ Avoid counterintuitive consequences of strong influence + indexicality?
- ▶ Consider how to merge (as quickly as possible) knowledge and/or belief within a social network.

## More further research

### Related work

- ▶ Use a less complex one-dimensional logic (without the plausibility dimension) to model properties distribution change within a network structure: current work with Jens U. Hansen (LORI-IV paper).
- ▶ Use a less complex one-dimensional logic (without the social network dimension): current work with Alexandru Baltag and Rasmus Rendsvig.
- ▶ Try to get an overview of logics for social science: future work with Sonja Smets.
- ▶ Consider how the network structure constrains the dynamics (example: odd circle forcing stabilization): future work with Johan van Benthem.
- ▶ Show how agents can come to know the structure of the network they are part of: future work with Nina Gierasimczuk.

Thank you 

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## References



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Talking your way into agreement: Belief merge by persuasive communication.

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